

# Reliable Distributed Consensus for Low-Power Multi-Hop Networks

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Second Marker

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# Overview

## Background

Wireless Sensor Networks  
Synchronous Transmissions  
Agreement Protocols  
Consensus

## Problems

Latency  
Reliability  
Configurability  
Replicability

## Contributions

Hybrid  
WISP  
WIMP

Background

# Wireless Sensor Network Nodes



One **Microcontroller** (MCU)

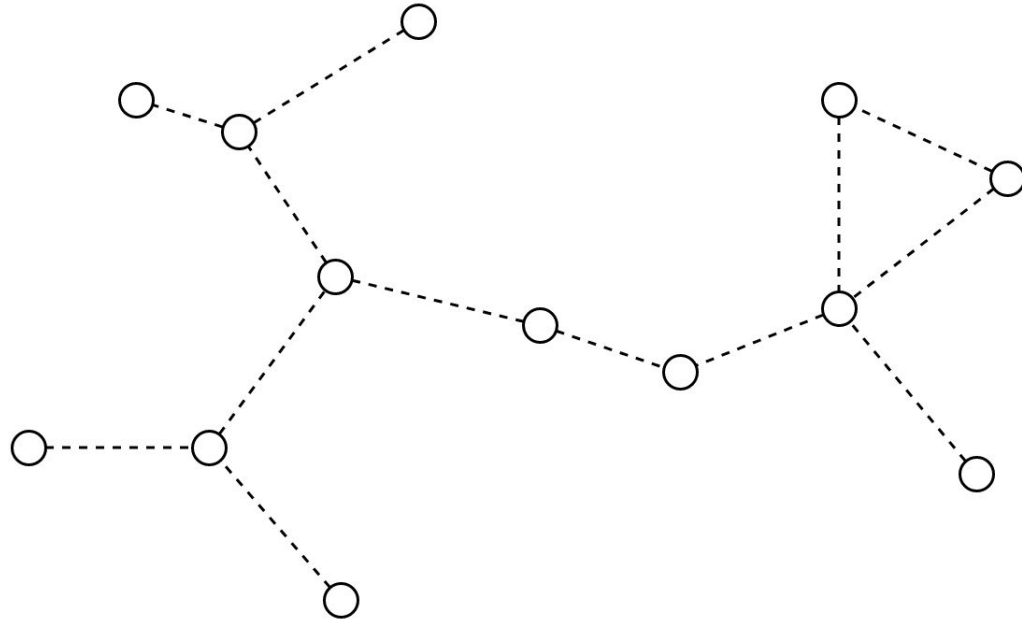
Multiple **Sensors/Actuators**

One Wireless Communicator or **Radio**

**Battery**

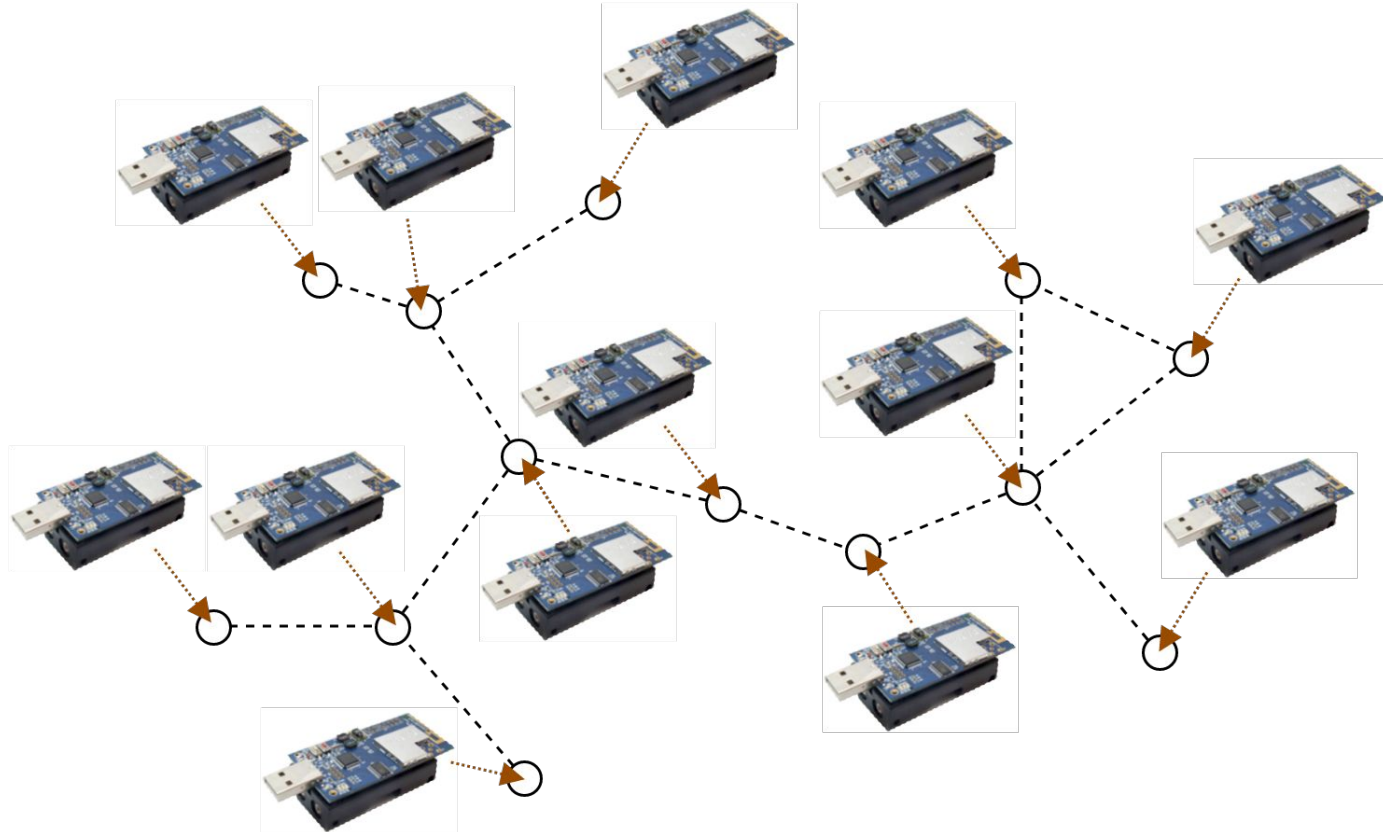
Background

# Wireless Sensor Networks



Background

# Wireless Sensor Networks



# Wireless Sensor Networks

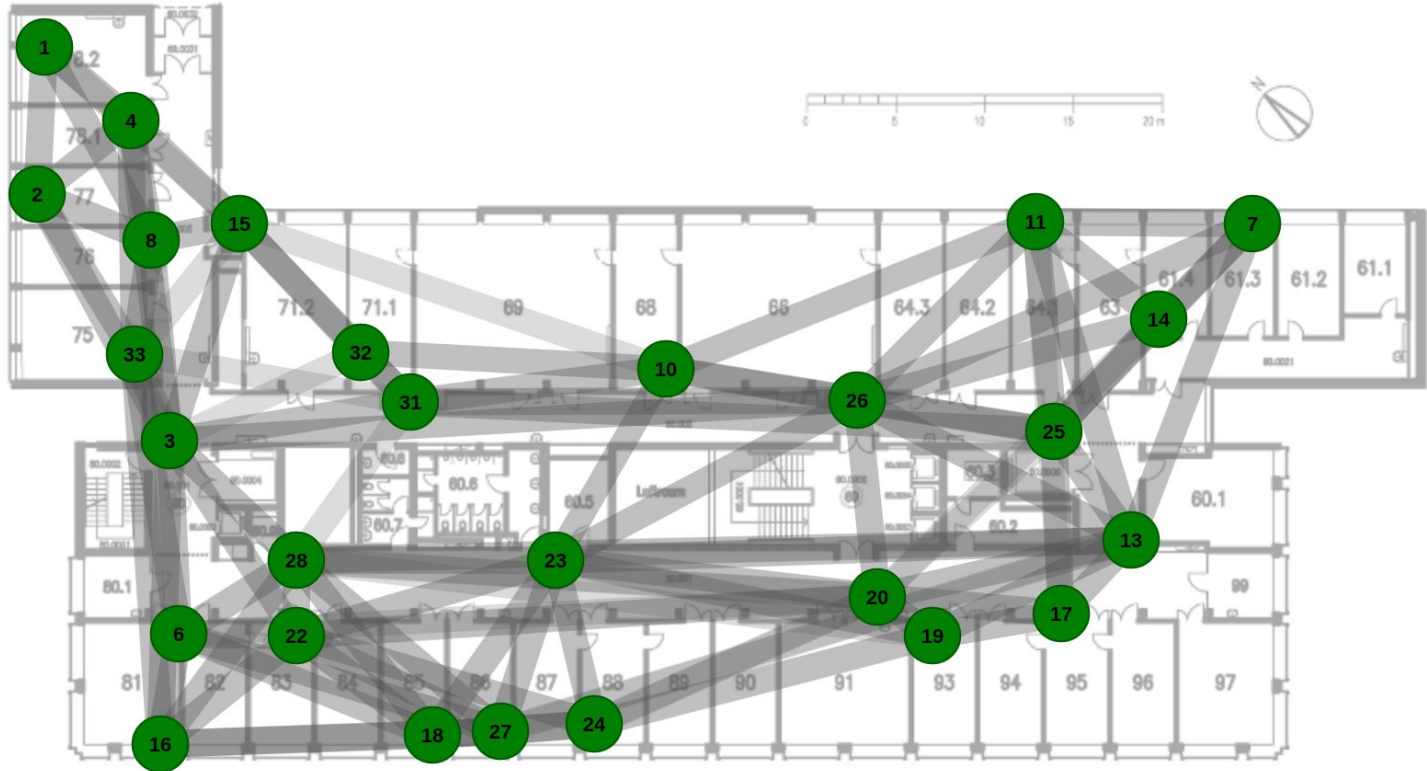
Use **low-power** radios

**Environment** causes **high interference**

Links are **unreliable**

## Background

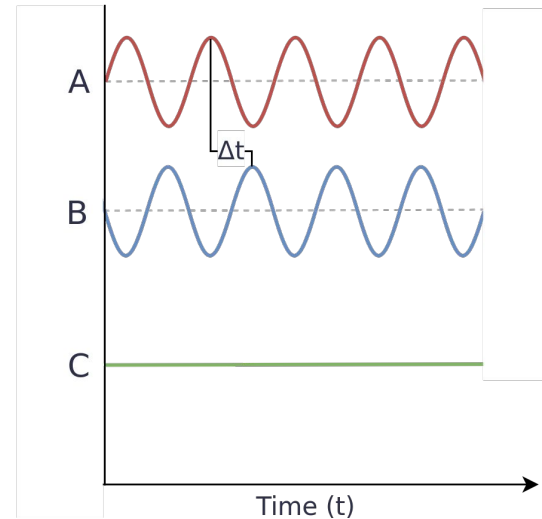
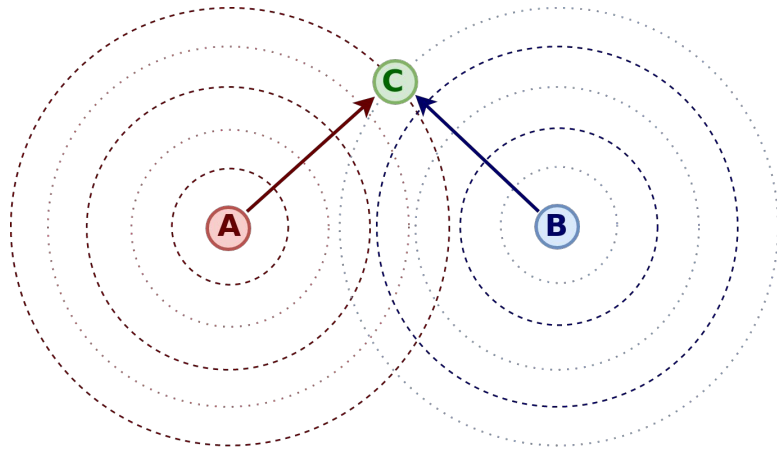
# FlockLab: a WSN testbed



# Broadcast Interference

**Phase difference** is perceived as a **time offset**  $\Delta t$

**Destructive Interference** for  $\Delta t$  odd multiple of  $\pi$

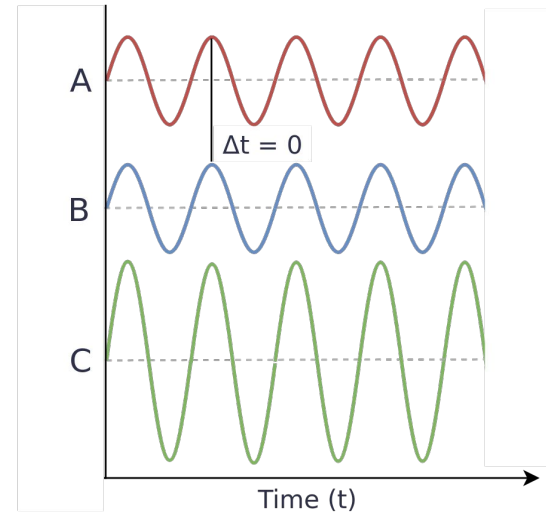
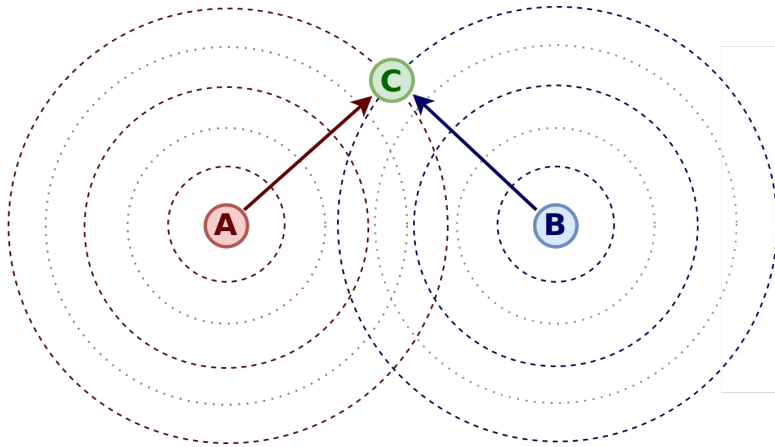




# Broadcast Interference

**Increase** the **reliability** of transmissions

**Constructive Interference** for  $\Delta t$  multiple of  $2\pi$



Background

# Synchronous Transmissions

**Constructive Interference**

+

**Capture Effect**

Repeatedly **flood packets** to the whole network

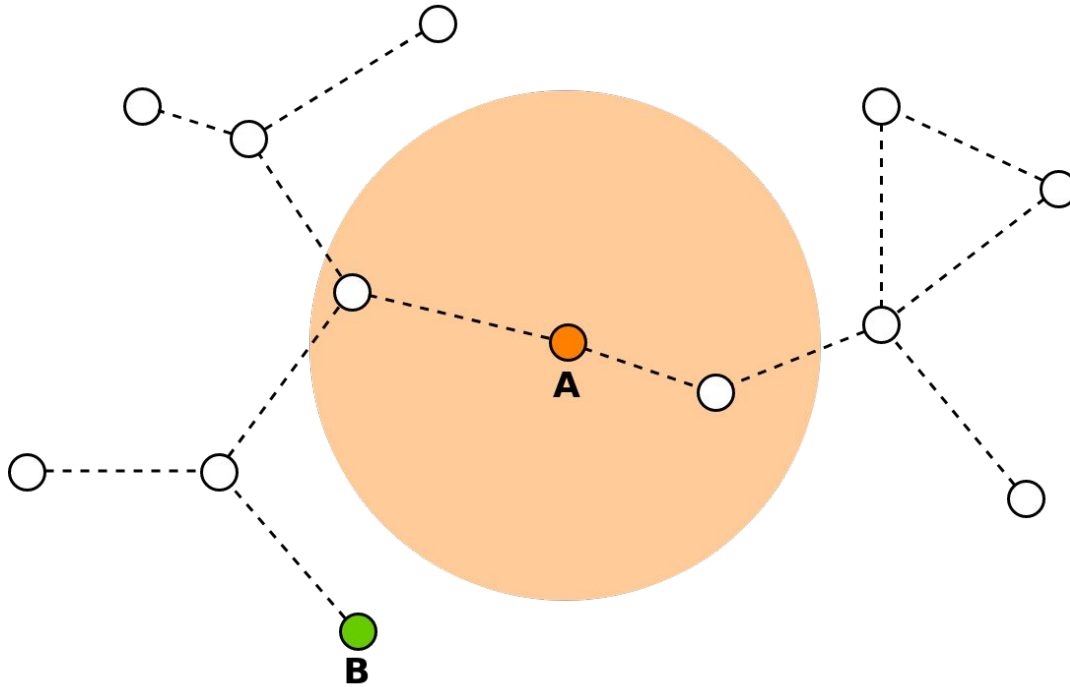
Enables **many-to-all communication**

**Route-less** packet broadcasts

Constant **winner at EWSN dependability competitions**

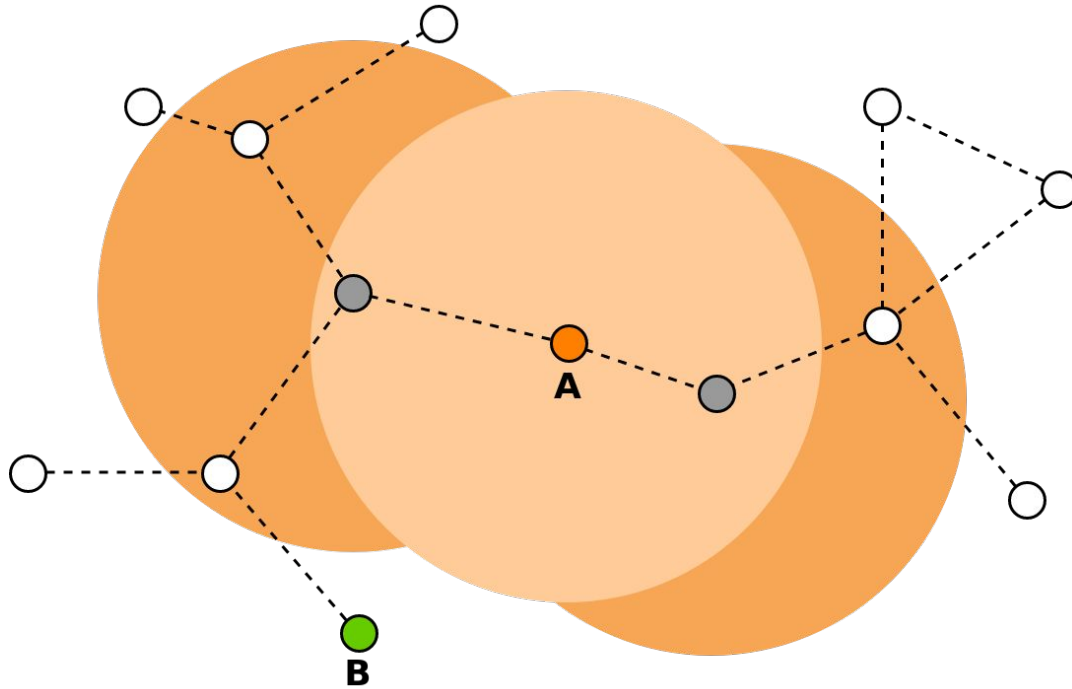
Background

# Synchronous Transmissions



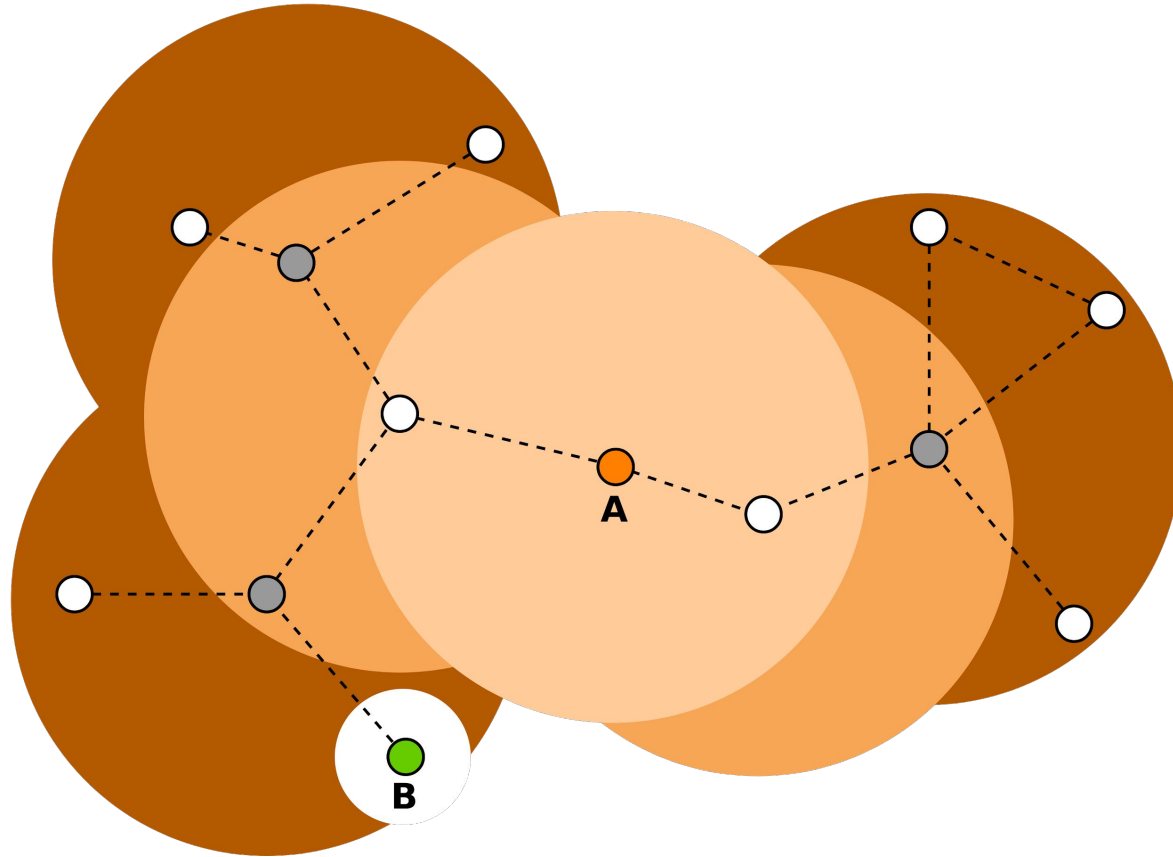
Background

# Synchronous Transmissions



Background

# Synchronous Transmissions



# Synchronous Transmission Primitives

**2011**

**Glossy**

**2013**

**Chaos**

2013

Splash

2013

SCIF

2014

P3

2015

Pando

2016

RedFixHop

2017

LiM

2017

Robust Flooding

2018

Mixer

2018

Codecast

Background

# Glossy

*F. Ferrari et al. (2011)*

**Synchronous Transmission (ST) primitive**

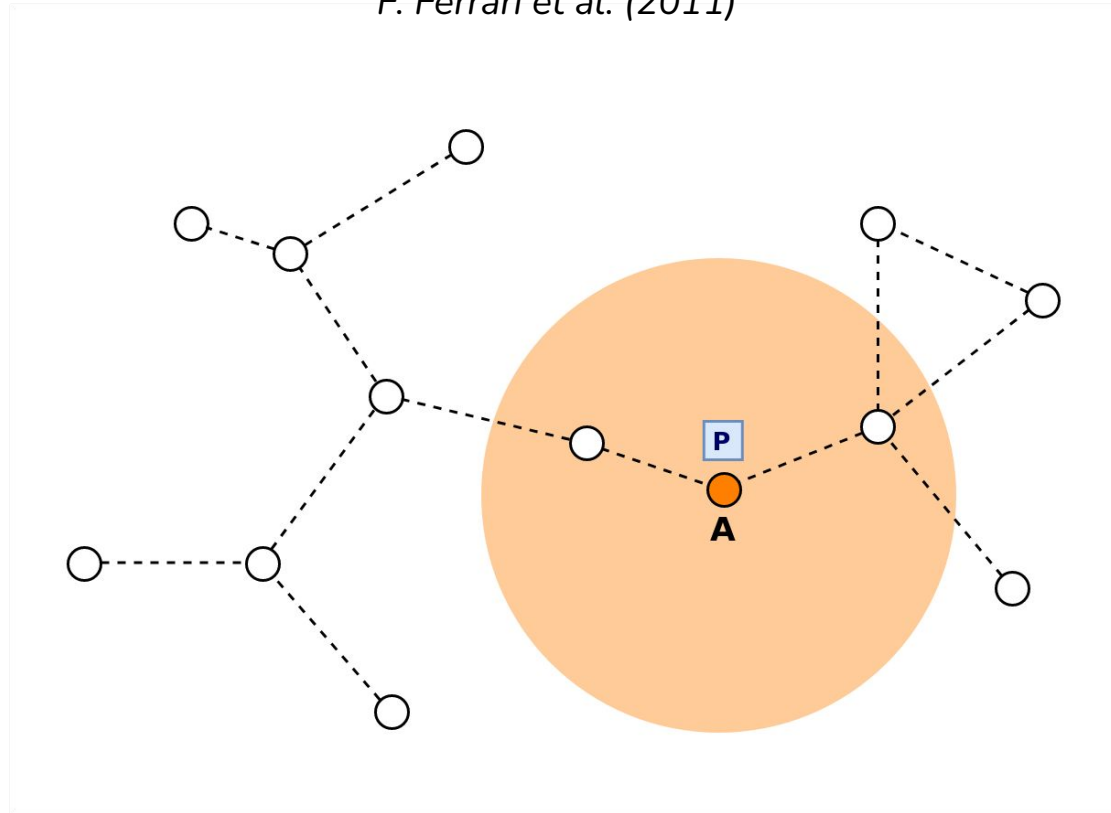
**One-to-all** data dissemination

**High reliability** guarantees

Background

# Glossy

*F. Ferrari et al. (2011)*

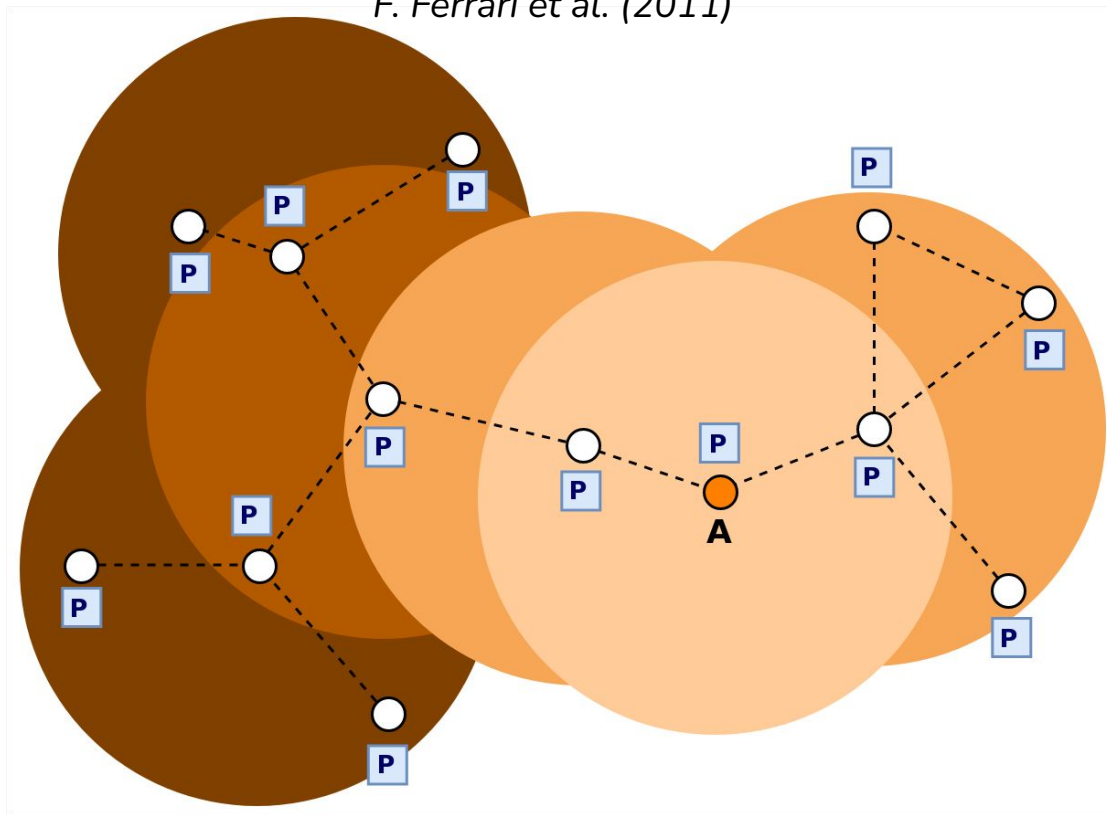




Background

# Glossy

*F. Ferrari et al. (2011)*



Problems

# Glossy

*F. Ferrari et al. (2011)*

**High latency** for all-to-all communication

**High energy** cost for dissemination of individual payloads

Background

# Chaos

*O. Landsiedel et al. (2013)*

**Many-to-all** data **dissemination**

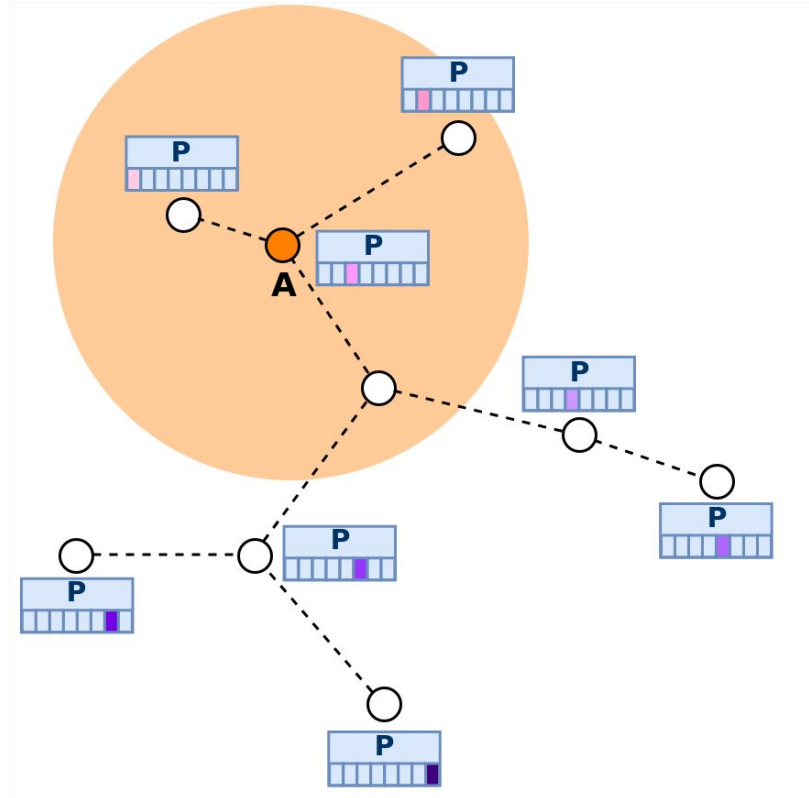
All **nodes contribute to the payload** being disseminated

**Power** and **time-efficient** flooding

## Background

# Chaos

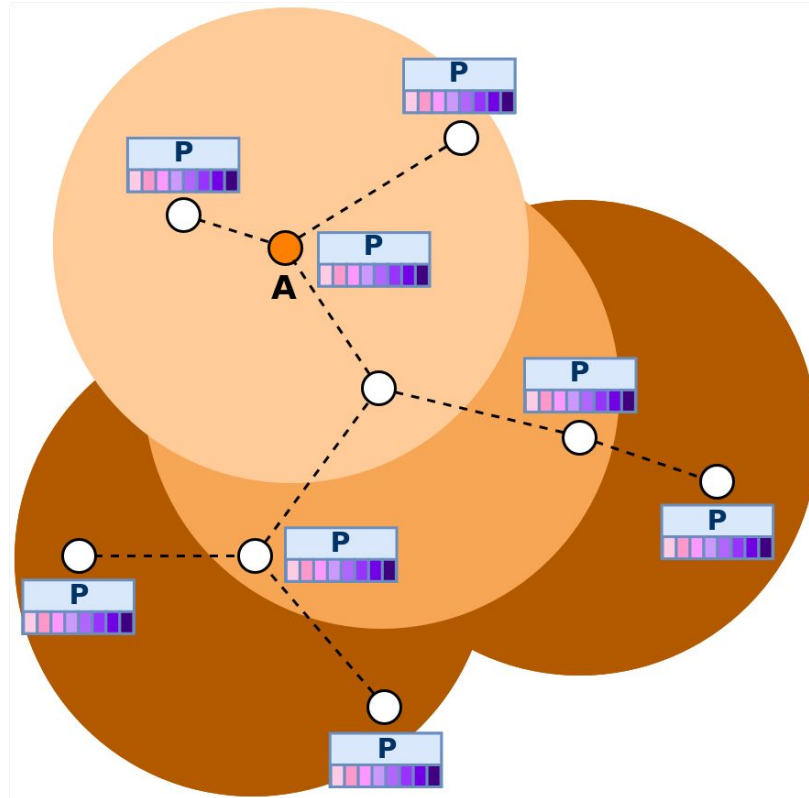
*O. Landsiedel et al. (2013)*



Background

# Chaos

*O. Landsiedel et al. (2013)*



Problems

# Chaos

*O. Landsiedel et al. (2013)*

**Unreliable** for one-to-all communication

**Unknown termination** time

# Voting Protocols

A **coordinator proposes a value** to the network

A network of nodes (**cohort**) **votes** on the proposed value

**Nodes conditionally commit**

**Atomic Commit** Protocols:

2PC & 3PC

**Consensus** Protocols:

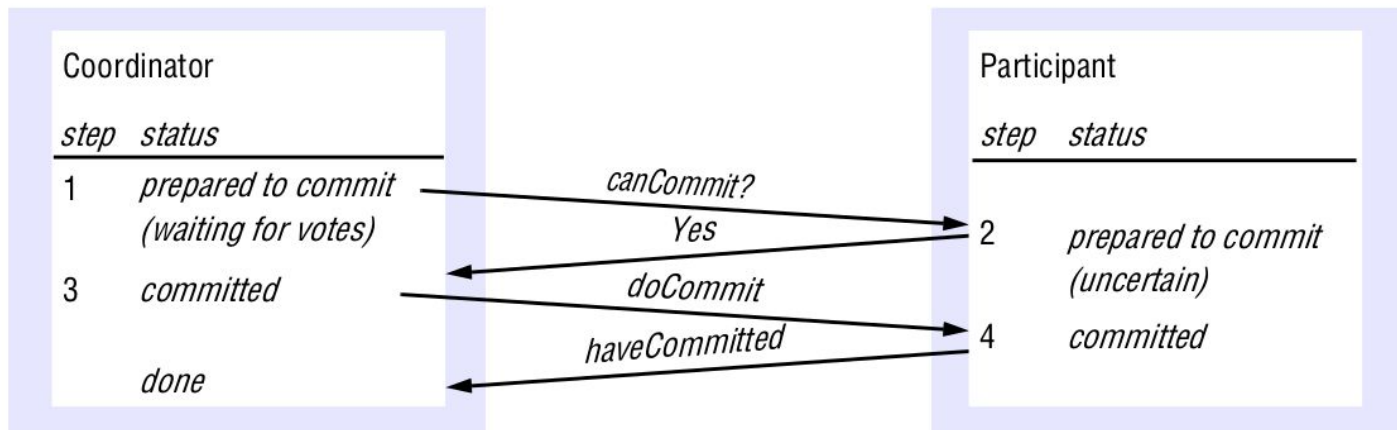
Paxos

# Two Phase Commit

J. Gray (1978)

All nodes will **commit** the **same value**

**Blocking protocol:** not guaranteed to terminate



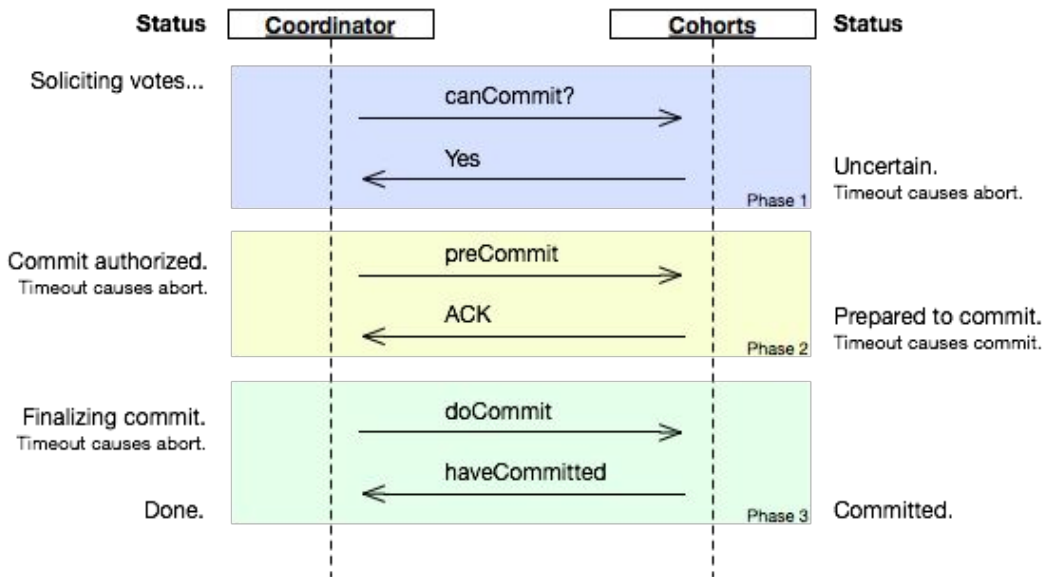


# Three Phase Commit

*D. Skeen (1981)*

Network might be  
**inconsistent**  
(**safety property**)

**Non-blocking:**  
guaranteed termination  
(**liveness property**)



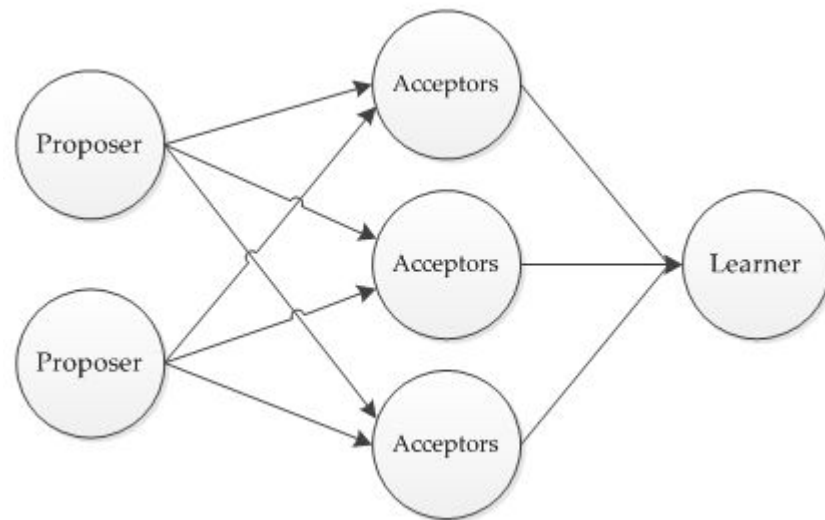
# Paxos

*L. Lamport (1998)*

**Values** are proposed to  
acceptors

Only a **quorum** of  
acceptors **must commit**

Committed **values** are  
**sent to learners**



# Voting Protocols

WSNs are **distributed asynchronous systems**

Common for broadcast **packets** to be **lost**

Nodes may become **unreachable with interference**

Background

# A<sup>2</sup>: Agreement in the Air

*B. Al Nahas et al. (2017) and V. Poirrot et al. (2019)*

Uses **Chaos ST primitive**

Implements **2PC and 3PC** protocols (2017)

Implements **consensus with WPaxos** (2019)

# A<sup>2</sup>: Agreement in the Air

*B. Al Nahas et al. (2017) and V. Poirot et al. (2019)*

**Chaos** ST primitive is **unreliable**

**No termination** guarantees

Greatly **disrupted by interference**

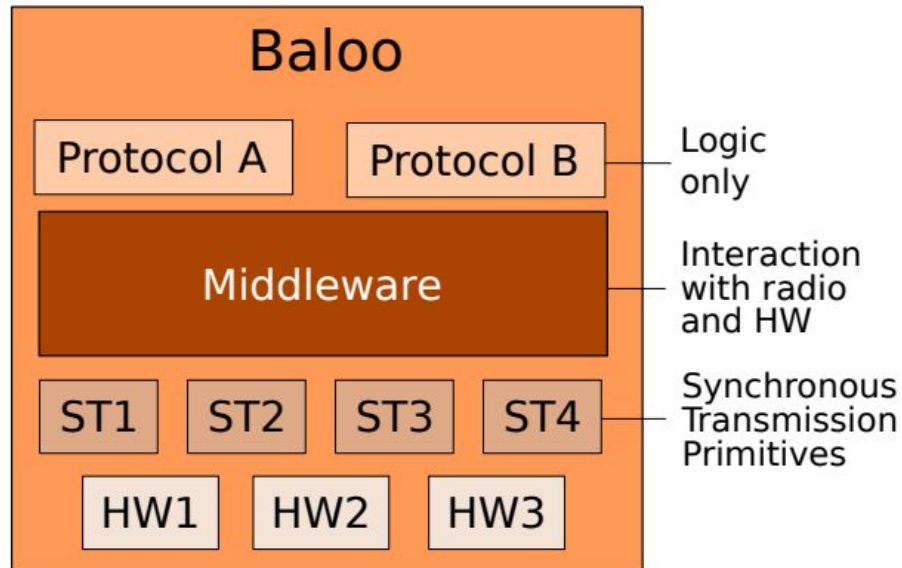
# Baloo

*R. Jacob et al. (2019)*

Proposes a **middleware** to **interact with radio** and hardware

Can be used with **multiple ST primitives**

**Protocols** implemented using **callback functions**



Problems

# Baloo

*R. Jacob et al. (2019)*

**Hard to use** same protocol with **more than one**  
**ST primitive**

Different **primitives** must be **scheduled differently**

**Protocols** must be in charge of **control and timing**

# Our Aim

- C1.** Protocols must be able to **easily switch between ST primitives**
- C2.** Create a **new ST primitive** which is able to provide the **robustness of Glossy** together with the **performance of Chaos**
- C3.** **Consensus** protocols must run **reliably on WSNs**
- C4.** Protocol reliability has to be **tested with replicable results**



# XPC: A Voting Protocol Coordinator

**Baloo requires a coordinator** for voting protocols

Handles flood **timing and** primitive **control**

Has **intermediate data representation** for primitives

**Addresses C1**

# XPC: A Voting Protocol Coordinator

**Baloo requires a coordinator** for voting protocols

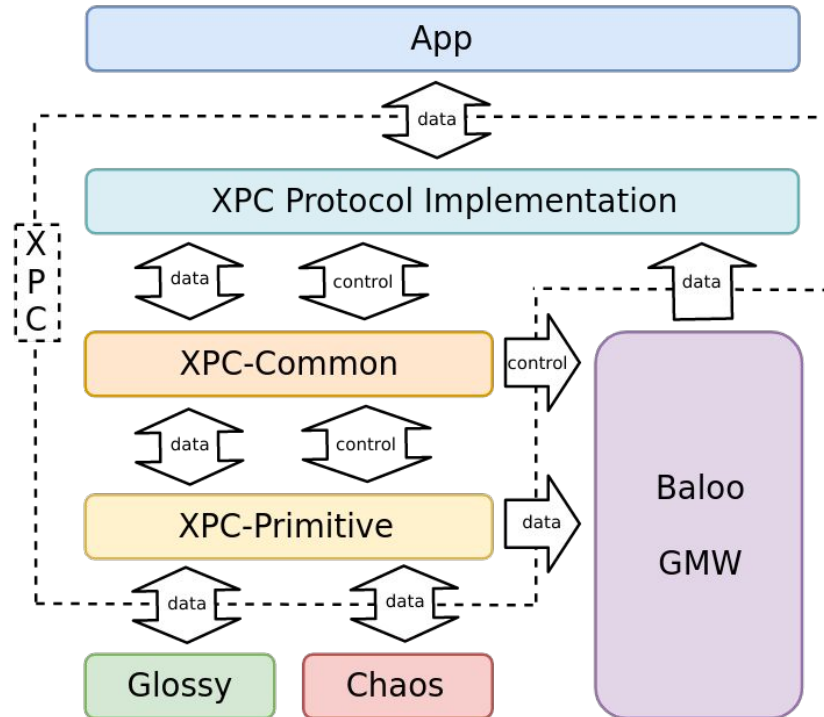
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Has **intermediate data representation** for primitives

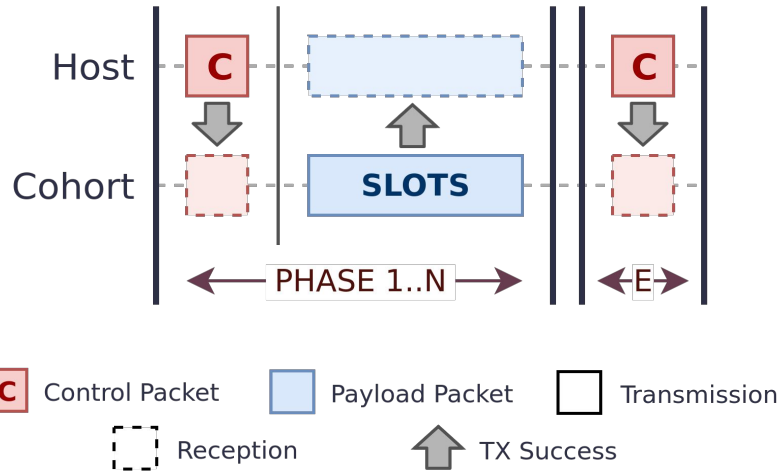
**Addresses C1**

“Protocols must be able to **easily switch between ST primitives**”

# XPC: A Voting Protocol Coordinator



# XPC Additions

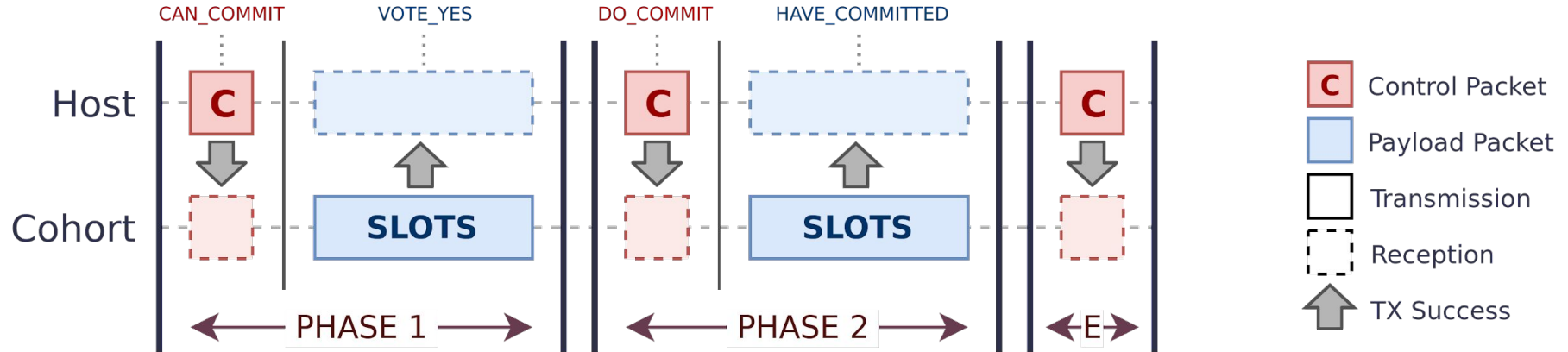


Single Initiator

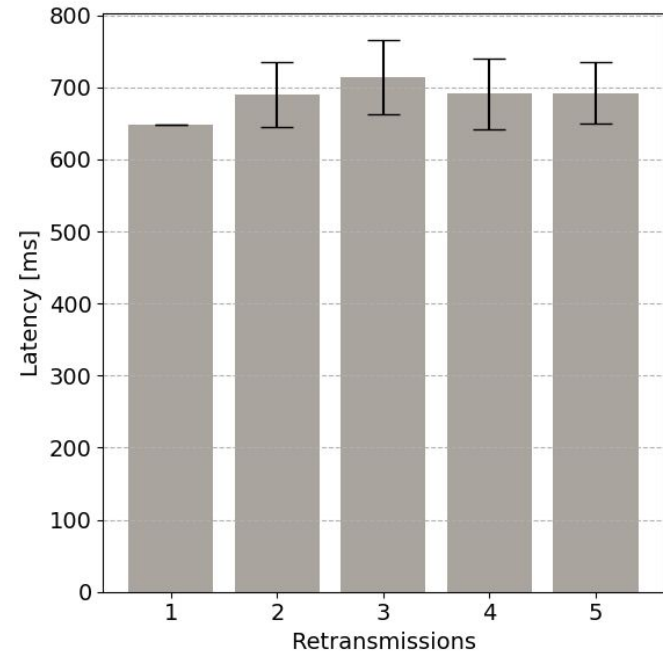
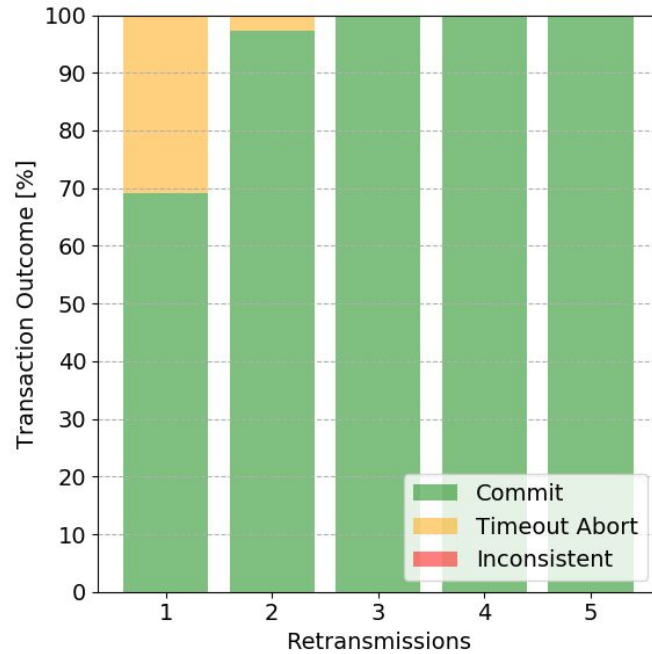
Retransmissions for Reliability

Additional **Final Round (E)**

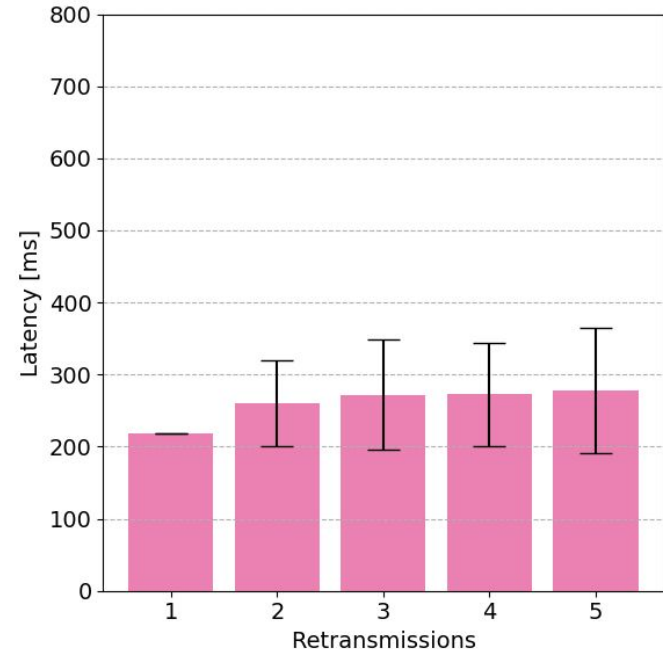
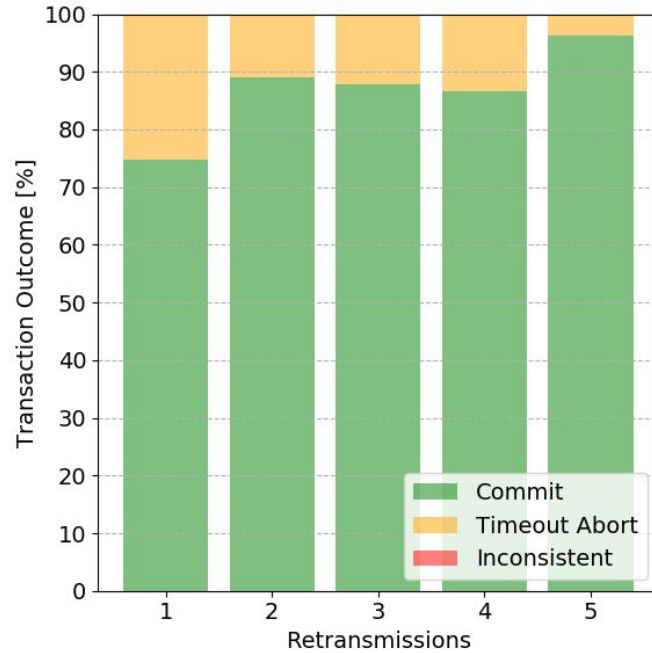
# Two Phase Commit with XPC



# 2PC-Glossy



# 2PC-Chaos

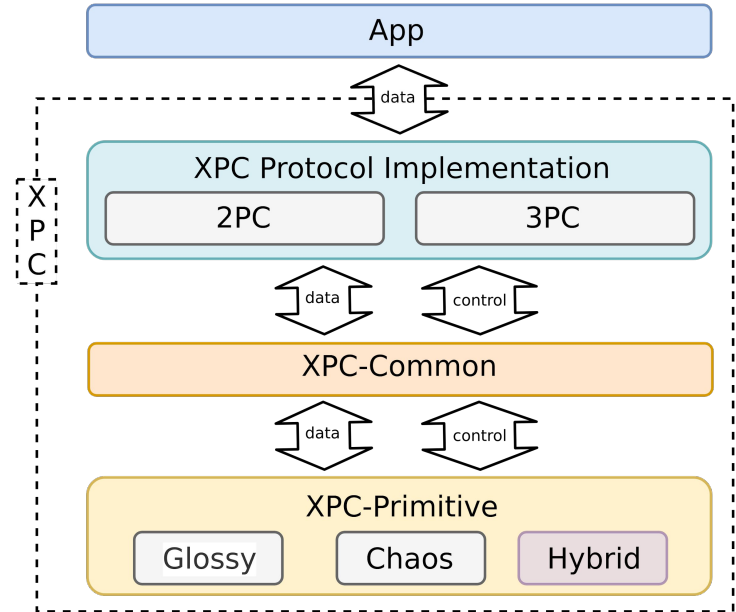


# Hybrid ST Primitive

Leverages **XPC** to schedule **first** a **Chaos** flood **and then Glossy** rounds

**Minimises latency** and **maximises reliability**

**Addresses C2**





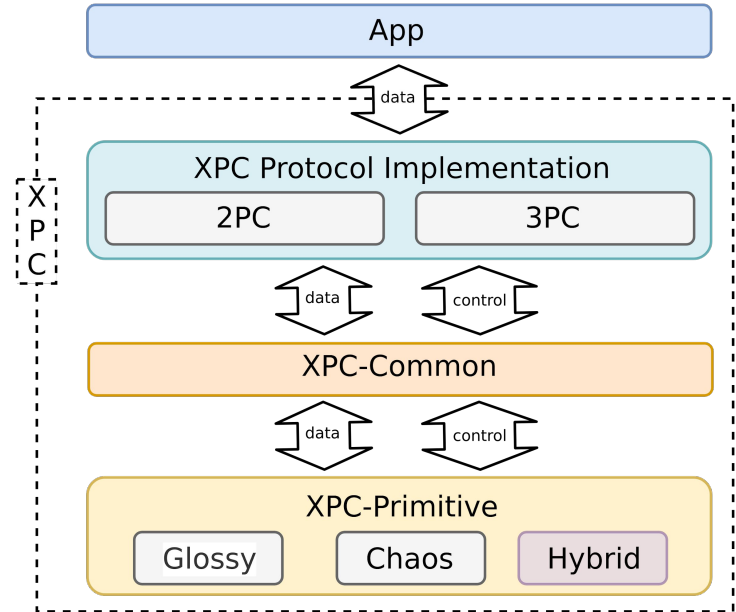
# Hybrid ST Primitive

Leverages **XPC** to schedule **first** a **Chaos** flood **and then Glossy** rounds

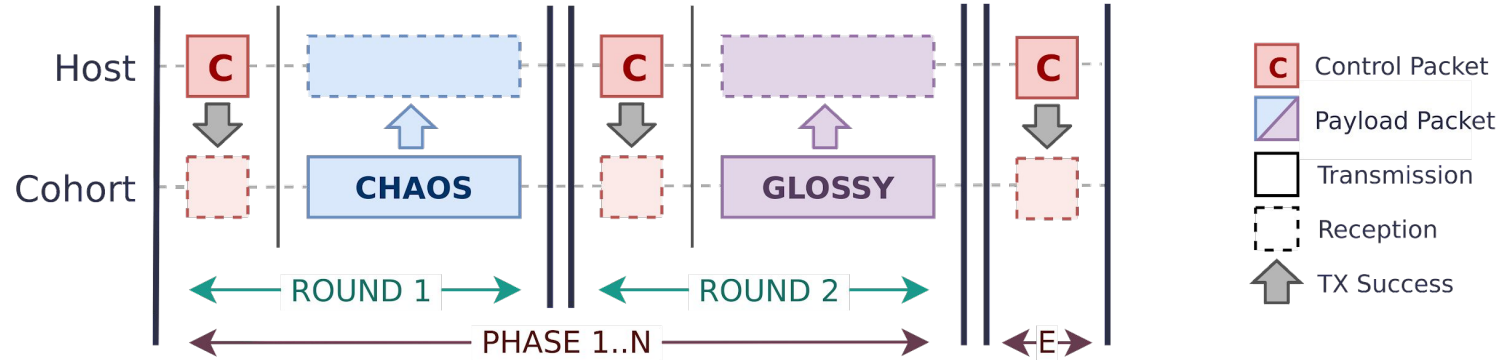
**Minimises latency** and **maximises reliability**

**Addresses C2**

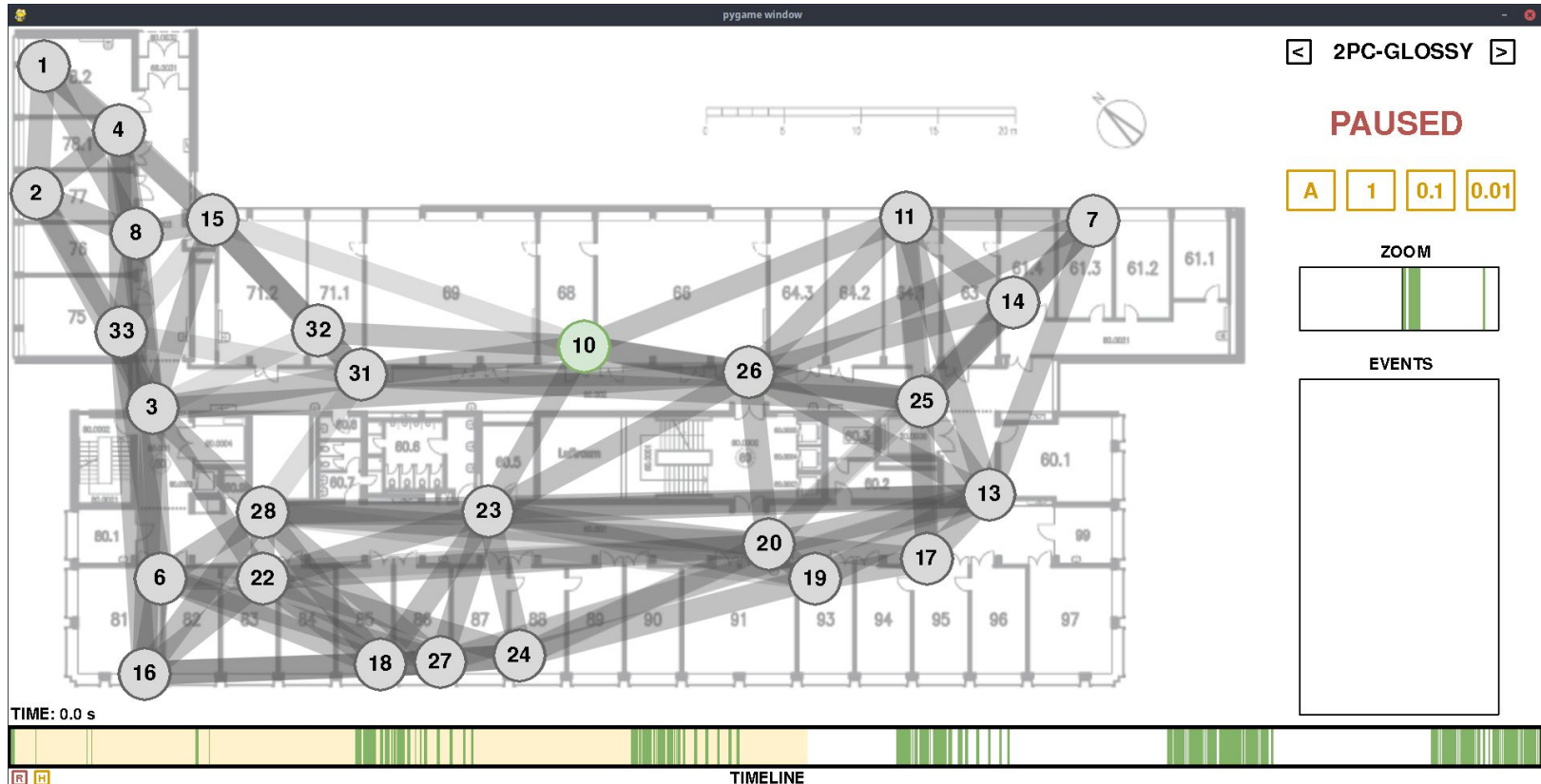
“Create a **new ST primitive** which is able to provide the **robustness of Glossy** together with the **performance of Chaos**”



# Hybrid ST Primitive



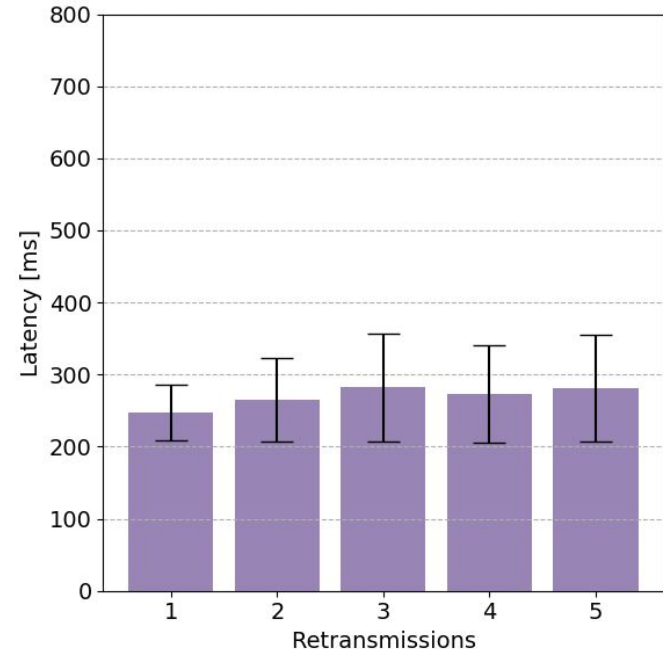
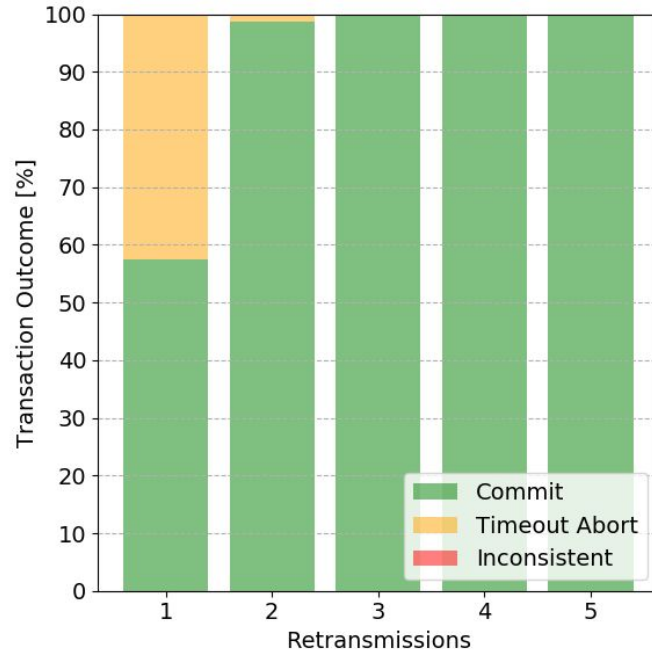
# Demo Interface: Flocklab Visualiser





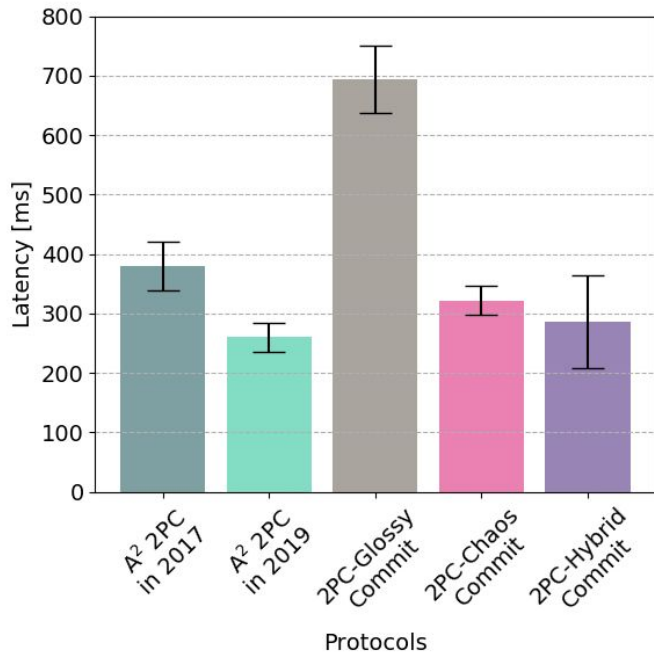
Contribution I

# 2PC-Hybrid

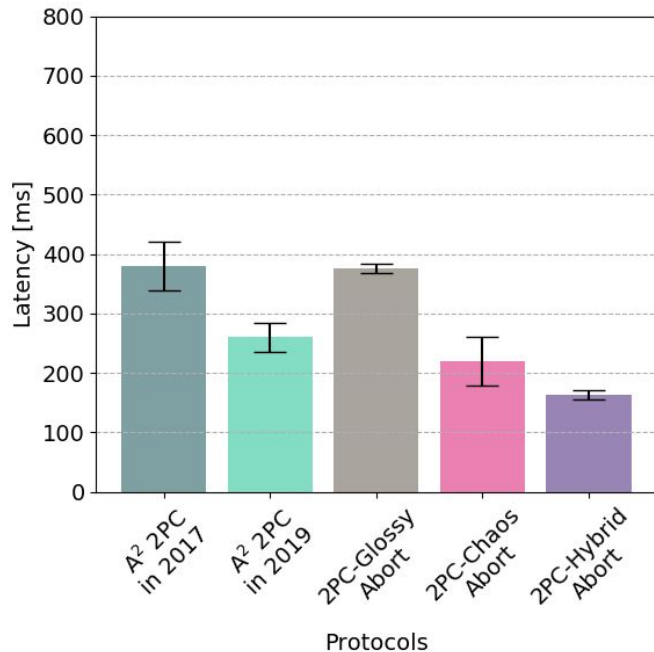


# Comparison with A<sup>2</sup>

## Transaction Commit



## Transaction Abort

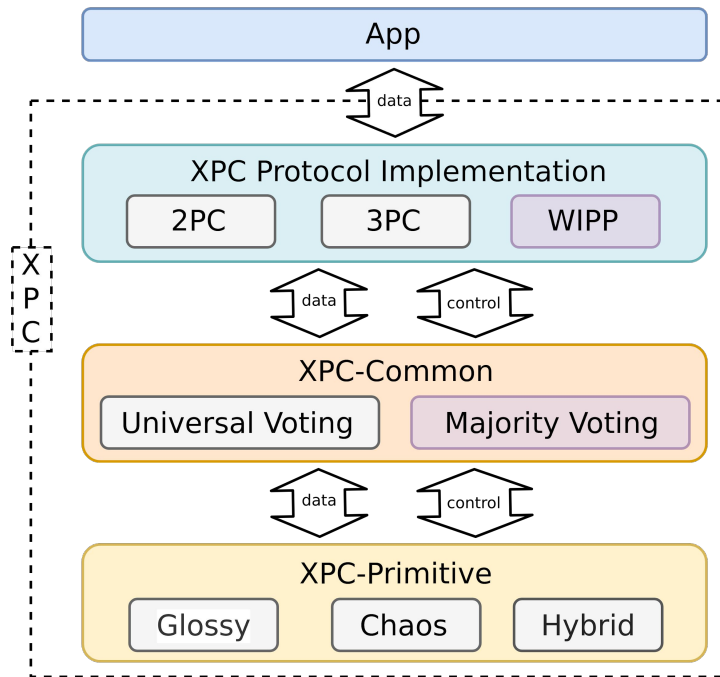


# WiPP: Wireless Part-time Parliament

Quorum-based **majority voting**

**Global dissemination** of  
committed values

Available as an **XPC Protocol**



# WISP: WiPP Simple Paxos

**WiPP** + **Hybrid** ST primitive

**Voting** phase **and** global **dissemination**

Satisfies **consensus properties**:

Validity

Integrity

Termination

Agreement

**Addresses C3**



# WISP: WiPP Simple Paxos

**WiPP** + **Hybrid** ST primitive

**Voting** phase **and** global **dissemination**

Satisfies **consensus properties**:

Validity

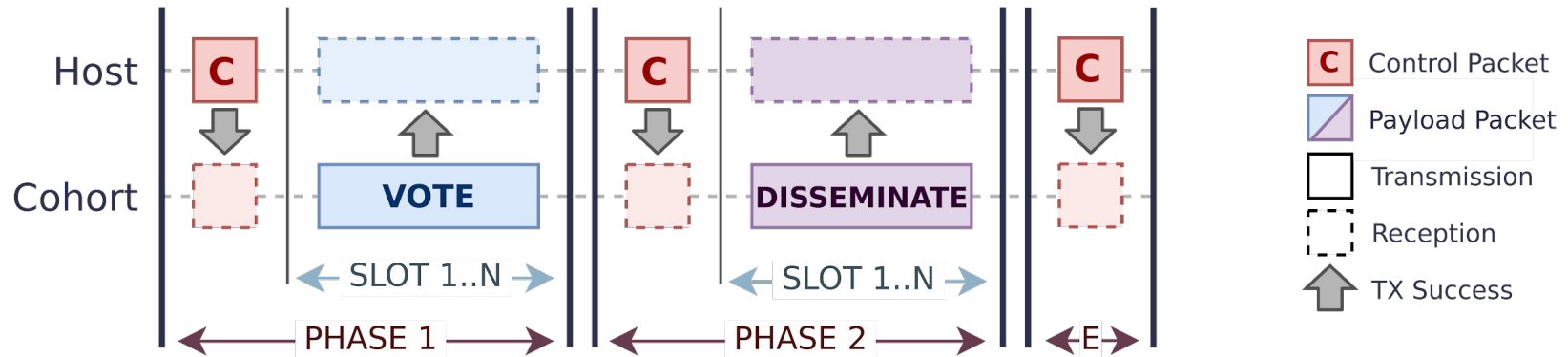
Integrity

Termination

Agreement

**Addresses C3**

“**Consensus** protocols must run **reliably on WSNs**”





# WISP Applications

Configuration management

Leader election

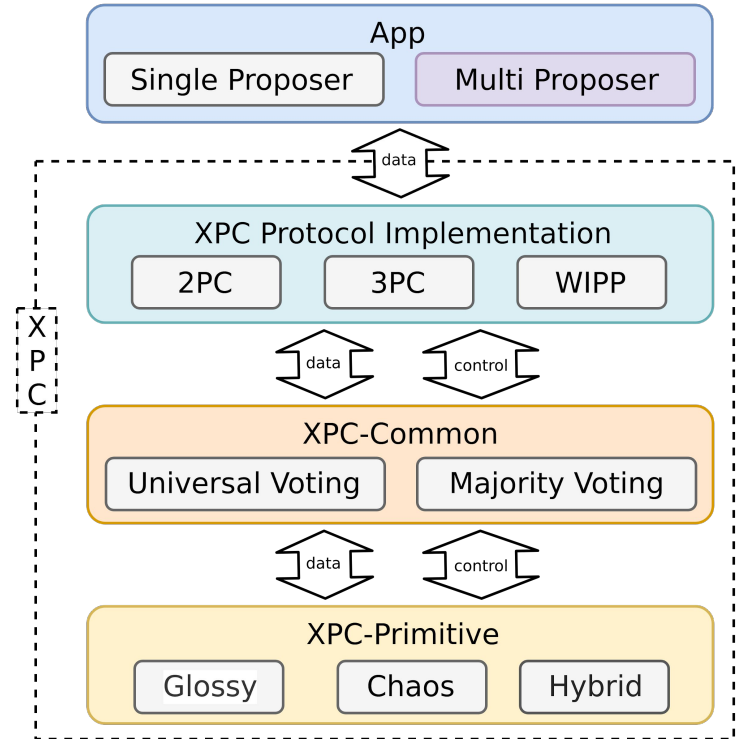
Local node clustering

Failure-free commits

# Multiple Proposers

Allows **proposals from any node**

Proposal use **contention slots**



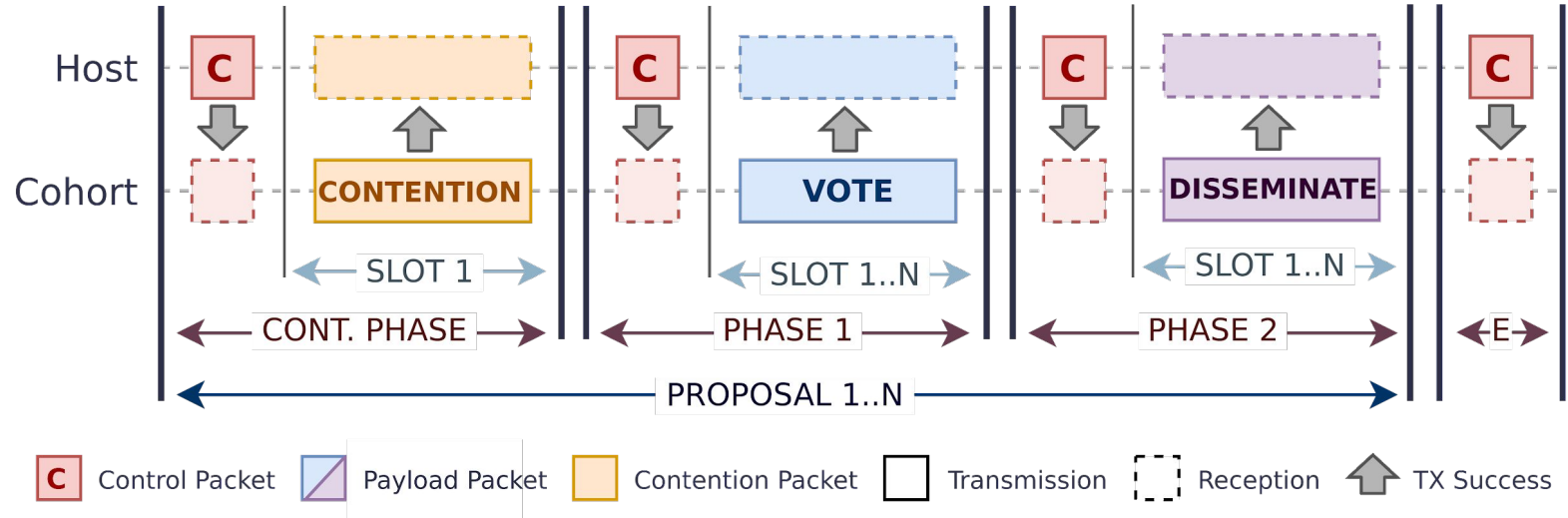
# WIMP: WiPP Multi Paxos

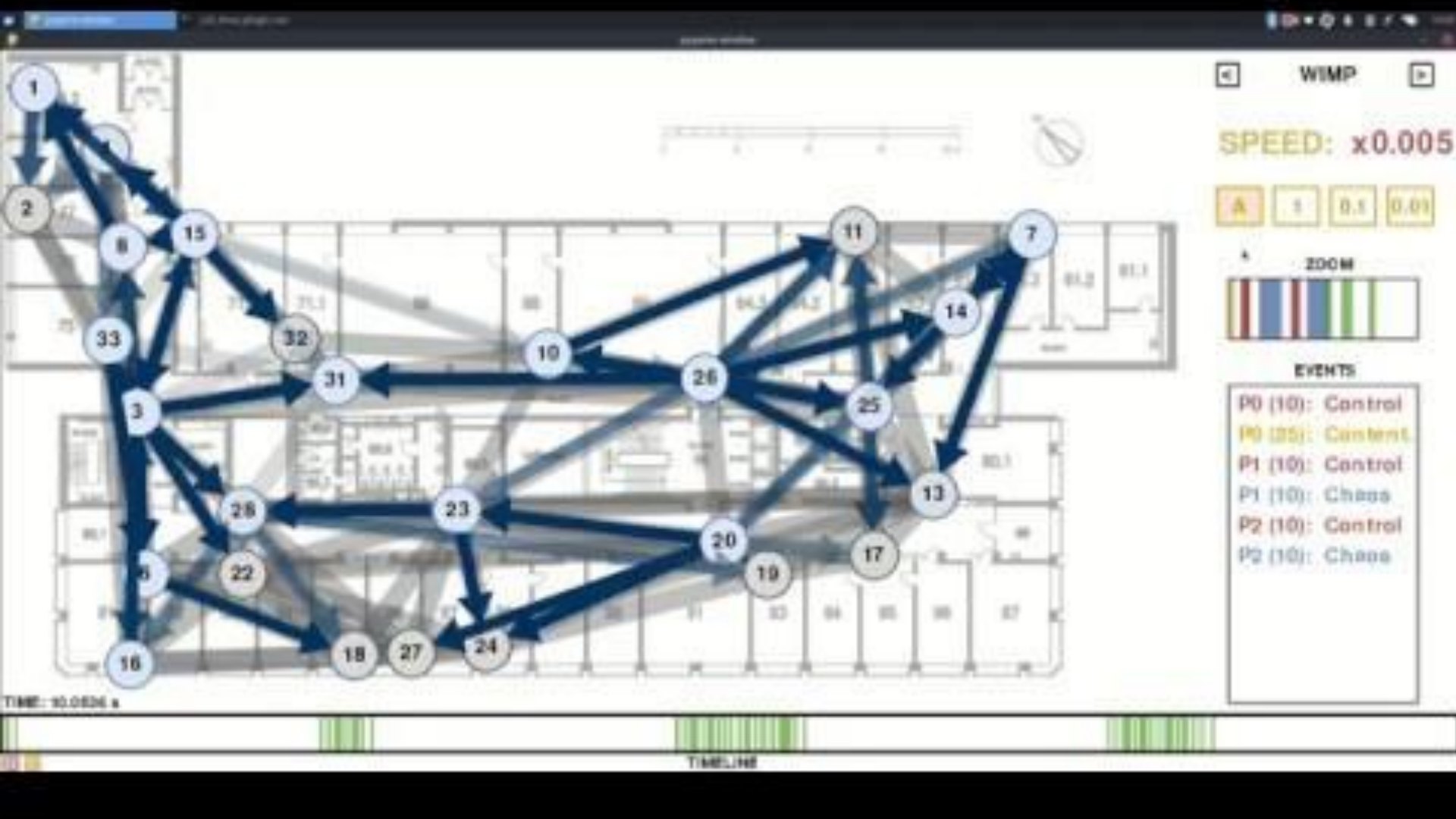
**Extension** of WISP

Allows **proposals** from **any network node**

Uses **global leader**

# WIMP: WiPP Multi Paxos

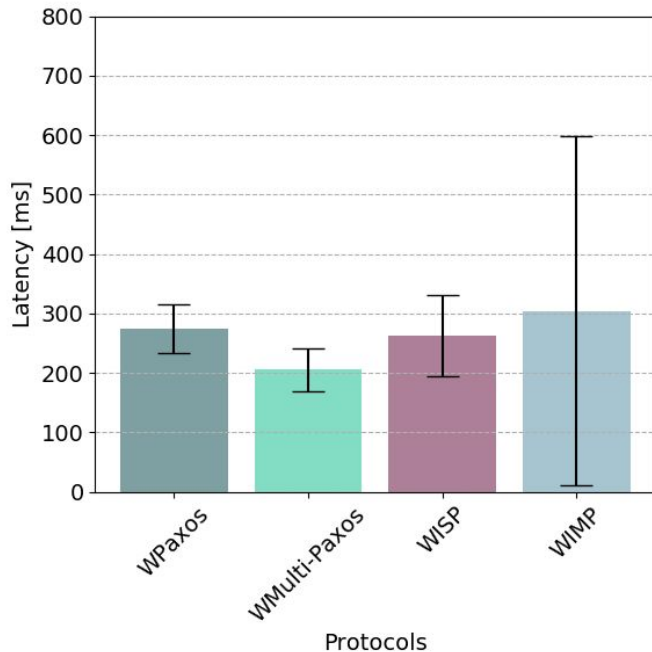




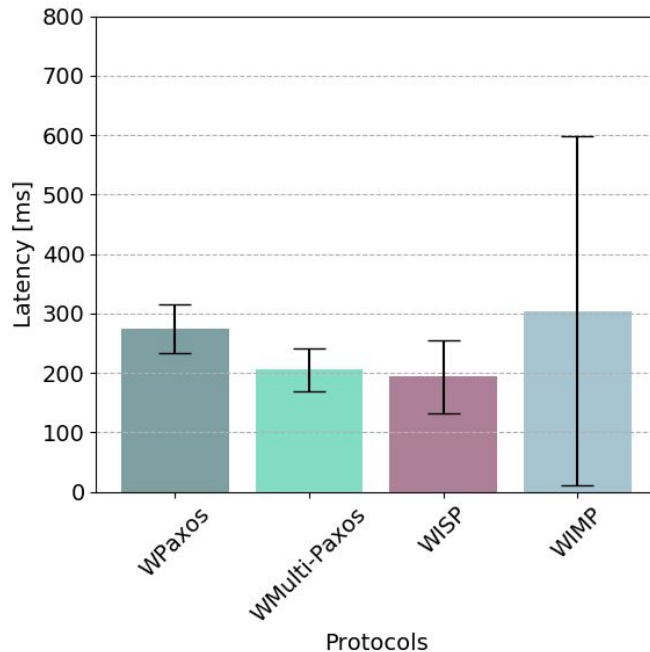


# Comparison with WPaxos ( $A^2$ )

## Transaction Commit



## Transaction Abort



Evaluation

# JamLab: Analysis with Interference

*C. A. Boano et al. (2011)*

To test **reliability** we **inject interference**

We use **JamLab**

Our protocols are **100% reliable with 1 interfering node**

**Addresses C4**

Evaluation

# JamLab: Analysis with Interference

*C. A. Boano et al. (2011)*

To test **reliability** we **inject interference**

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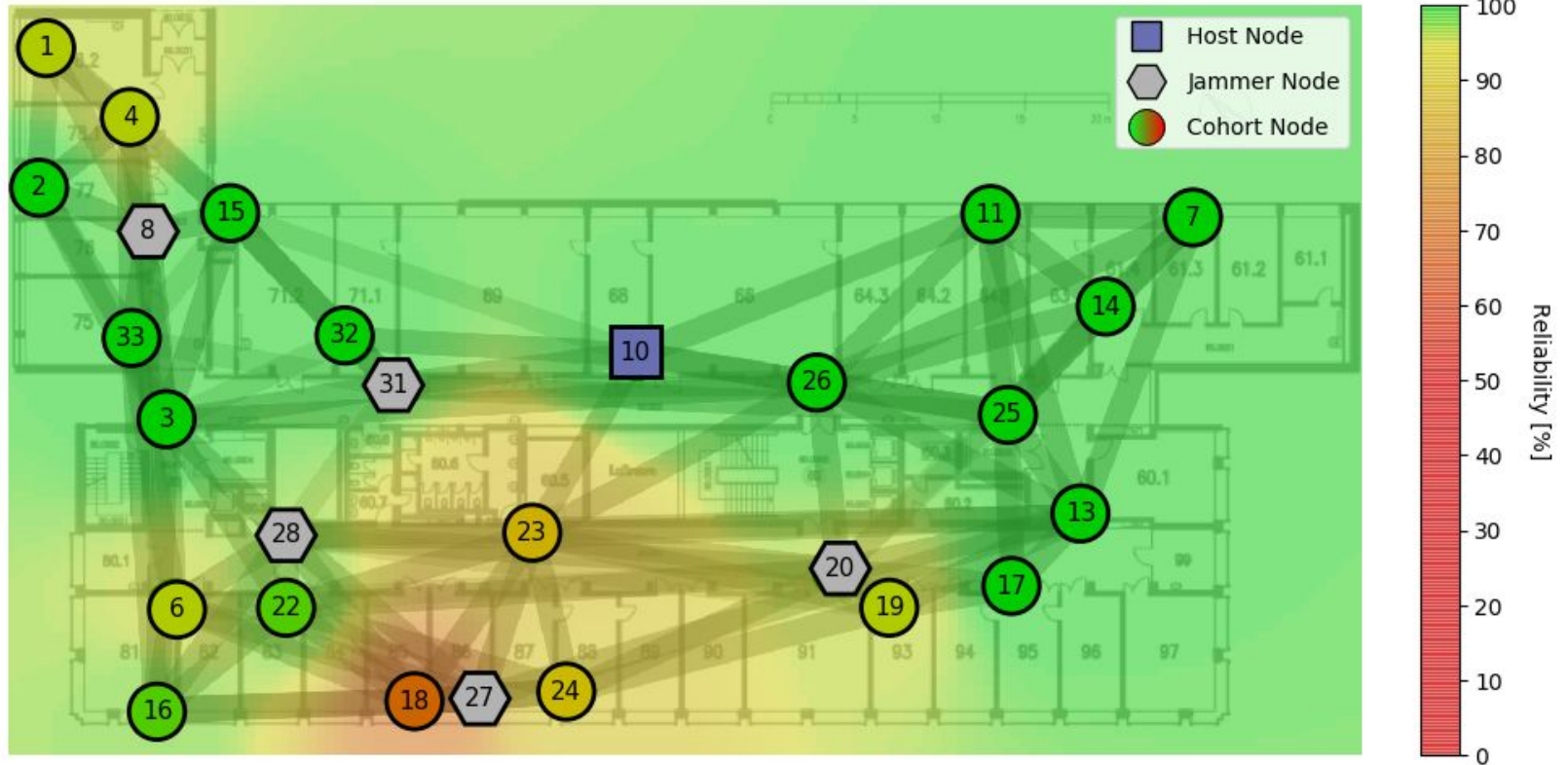
Our protocols are **100% reliable with 1 interfering node**

**Addresses C4**

“Protocol reliability has to be **tested with replicable results**”



# Multiple Interfering Nodes



Conclusion

# Our Contributions

**Hybrid**  
(and XPC)

**WISP**  
(and WiPP)

**WIMP**  
(and Multi-Proposer)

Conclusion

# Future Work

**Group membership**

**XPC** protocols **as a service**

Submit for **publication to IPSN**

Conclusion

Q & A



# Security Concerns

ST primitives are **not secure**

A lot of research in WSN Security:

C. Chu et al. (2010). Practical ID-based Encryption for Wireless Sensor Network.

R. Gustavo et al. (2012). Asymmetric Encryption in Wireless Sensor Networks.

D. Shubhangi et al. (2015). Security in Wireless Sensor Network Using Cryptographic Techniques.

M. Elhoseny et al. (2016). An energy efficient encryption method for secure dynamic WSN. Security and Communication Networks.

K. C. Hewage et al. (2017). Protecting Glossy-Based Wireless Networks from Packet Injection Attacks.

K. Tsai et al. (2018). A Light Weight Data Encryption Method for WSN Communication.

Additional

# Multi-Node Interference

## Termination

Eventually each correct process decides a value

## Agreement

All correct processes decide on the same value

## Integrity

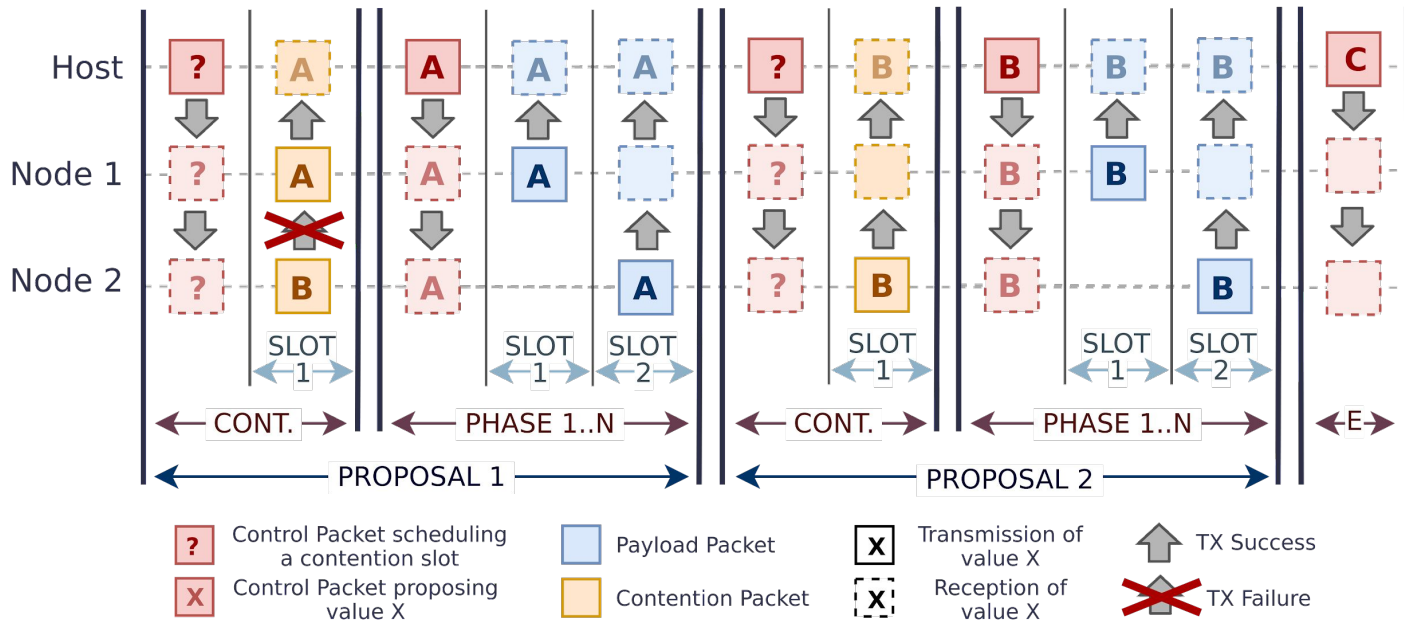
A process decides at most on one value

## Validity

If a process decides on a value, then it must have been proposed by some process

Additional

# Multiple-Proposers in Action



Additional

# Safety in Multiple-Proposers

**PPB:** Pending Proposal Bit

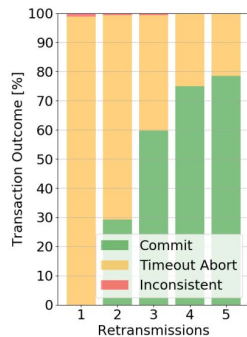
Used when Nodes wish to propose

**PVB:** Proposed Value Bit

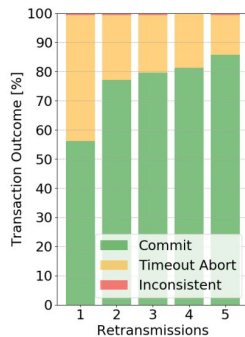
Used for proposal acknowledgement

Additional

# 2PC-Chaos Overview



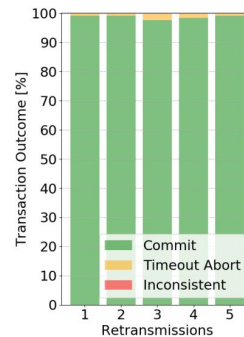
(a) 25ms slots



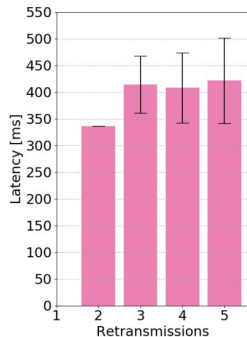
(b) 50ms slots



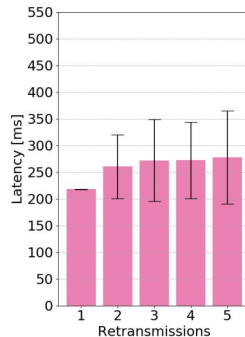
(c) 100ms slots



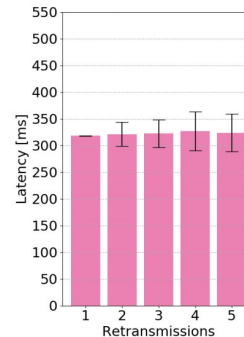
(d) 200ms slots



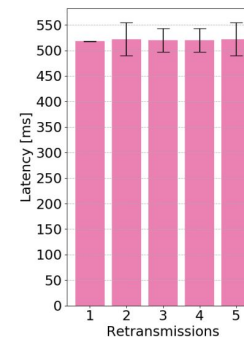
(a) 25ms slots



(b) 50ms slots



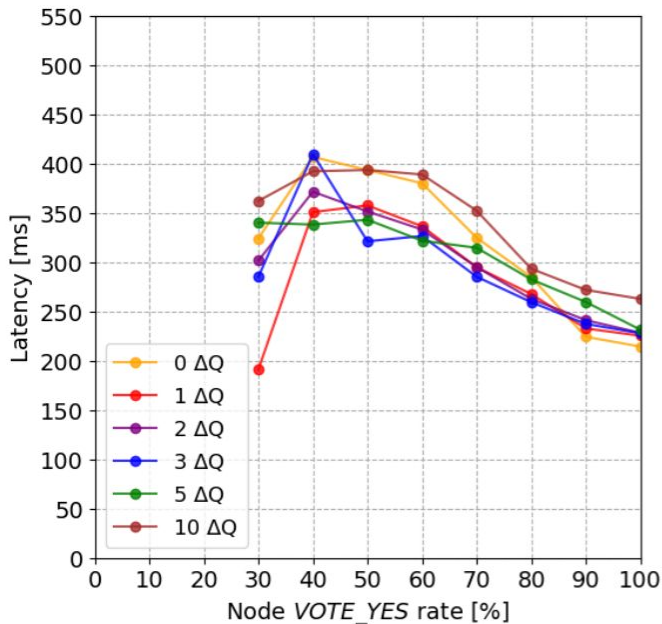
(c) 100ms slots



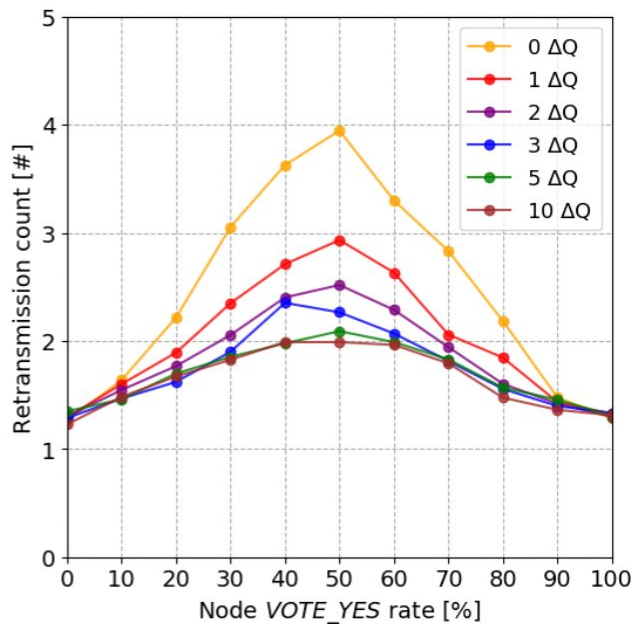
(d) 200ms slots

Additional

# Majority Voting Delta-Q



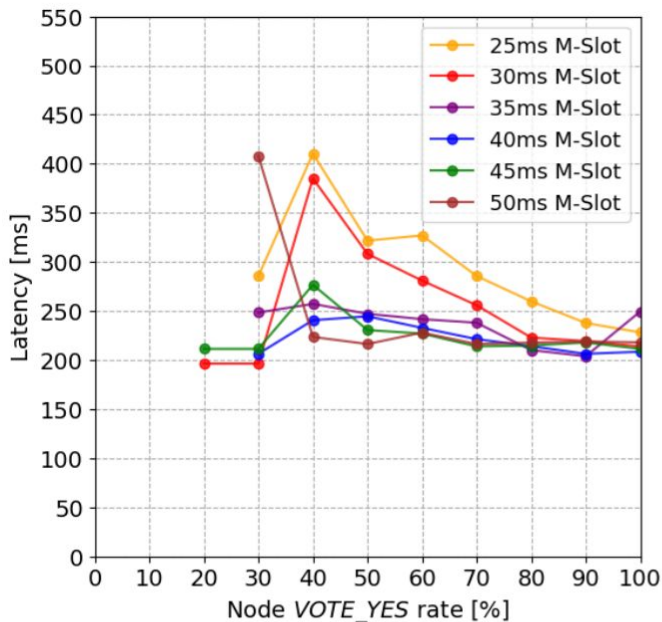
(a) Average commit latency



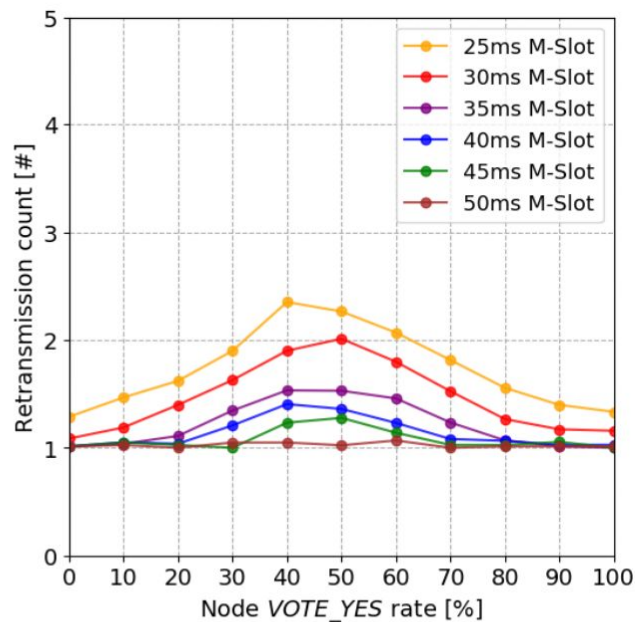
(b) Average number of retransmissions

Additional

# Majority Voting M-Slots



(a) Average commit latency



(b) Average number of retransmissions

# Interference Overview

Wifi Interference	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	100.00	383.07	P1: 92.19 P2: 90.54	P1: 2.29 P2: 2.19
3PC-Hybrid	100.00	606.20	P1: 91.32 P2: 92.92 P3: 89.58	P1: 2.26 P2: 1.96 P3: 2.31
WISP	100.00	270.93	MP: 87.79 DP: 91.98	MP: 1.05 DP: 1.44
WIMP	100.00	360.88	MP: 86.09 DP: 91.98	MP: 1.04 DP: 1.70
Microwave	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	100.00	396.06	P1: 92.97 P2: 91.82	P1: 2.07 P2: 2.17
3PC-Hybrid	100.00	586.67	P1: 92.09 P2: 91.58 P3: 92.30	P1: 1.95 P2: 2.12 P3: 1.95
WISP	100.00	342.53	MP: 78.90 DP: 91.14	MP: 1.15 DP: 2.23
WIMP	100.00	377.31	MP: 86.22 DP: 91.69	MP: 1.01 DP: 1.66

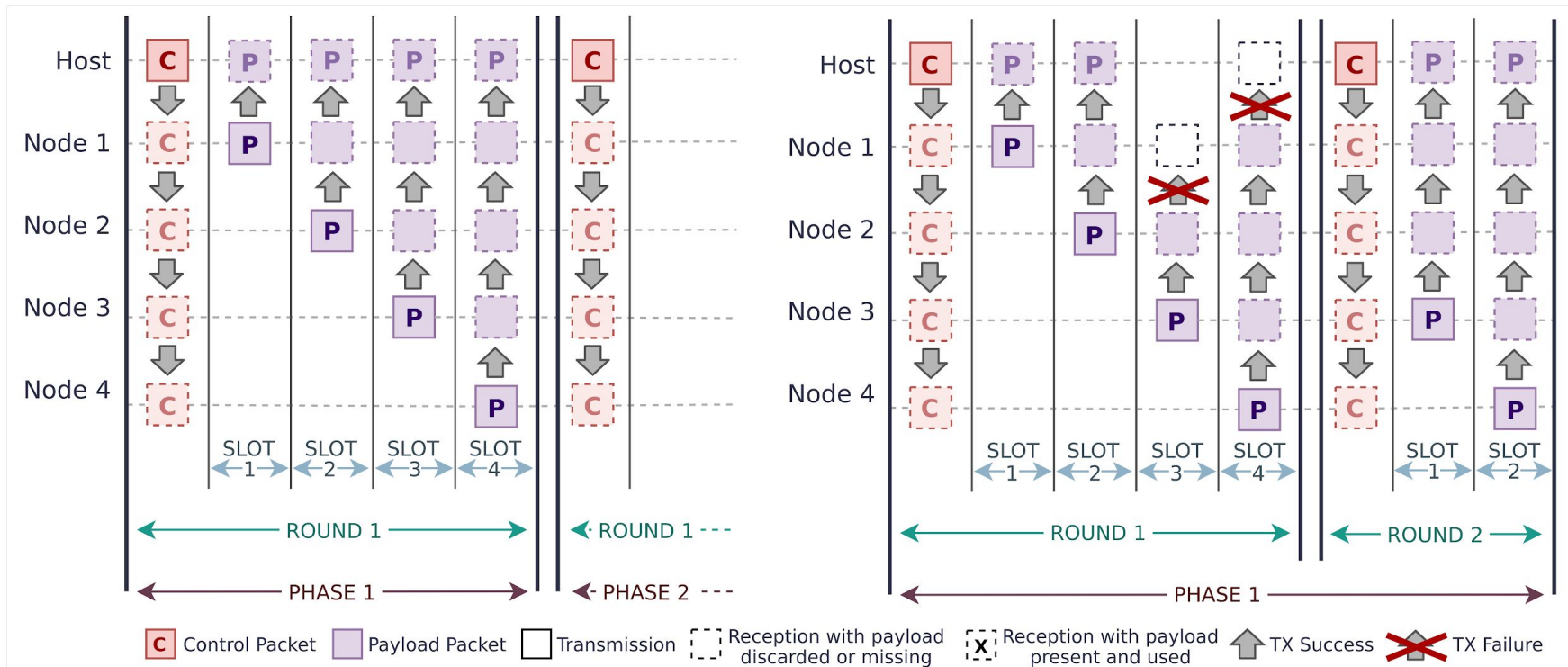


# Multi-Node Interference

Microwave (2 Nodes)	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	100.00	1027.18	P1: 82.91 P2: 81.54	P1: 4.64 P2: 5.04
WISP	98.33	625.37	MP: 56.49 DP: 82.85	MP: 1.57 DP: 4.52
Microwave (3 Nodes)	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	95.92	1025.86	P1: 80.26 P2: 81.43	P1: 4.95 P2: 5.11
WISP	92.86	756.73	MP: 47.00 DP: 81.37	MP: 1.97 DP: 5.22
Microwave (4 Nodes)	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	57.14	1221.30	P1: 79.59 P2: 79.92	P1: 6.72 P2: 6.11
WISP	59.18	1024.20	MP: 42.06 DP: 78.83	MP: 2.27 DP: 6.91
Microwave (5 Nodes)	Reliability (%)	Latency (ms)	Chaos Coverage (%)	Avg. Retr.
2PC-Hybrid	47.22	1291.68	P1: 75.51 P2: 79.37	P1: 7.85 P2: 7.50
WISP	59.09	1178.84	MP: 41.67 DP: 73.88	MP: 2.83 DP: 7.67

Additional

# XPC with Glossy





Additional

# XPC with Chaos

